

ShaperProbe: End-to-end Detection of ISP Traffic Shaping using Active Methods*

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ABSTRACT

We present an end-to-end measurement method for the detection of traffic shaping. Traffic shaping is typically implemented using token buckets, allowing a maximum burst of traffic to be serviced at the peak capacity of the link, while any remaining traffic is serviced at a lower shaping rate. The contribution of this paper is twofold. First, we develop an active end-to-end detection mechanism, referred to as ShaperProbe, that can infer whether a particular path is subject to traffic shaping, and in that case, estimate the shaper characteristics. Second, we analyze results from a large-scale deployment of ShaperProbe on M-Lab over the last 24 months, detecting traffic shaping in several major ISPs. Our deployment has received more than one million runs so far from 5,700 ISPs.

Categories and Subject Descriptors

C.2.3 [Computer-Communication Networks]: Network Operations—*Network monitoring*

General Terms

Measurement, Performance

Keywords

Active probing, traffic shaping, inference

1. INTRODUCTION

The increasing penetration of broadband access technologies, such as DSL, DOCSIS and WiMAX, provides users with a wide range of upstream and downstream service rates. Broadband users need to know whether they actually get the service rates they pay for. On the other hand, ISPs now have an extensive toolbox of traffic management mechanisms they can apply to their customers' traffic: application classifiers, schedulers, active queue managers etc. In this paper

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we focus on a class of such mechanisms referred to as *traffic shapers* or *traffic policers*.¹

A traffic shaper is a single-input single-output packet forwarding module that behaves as follows: Consider a link of capacity C bps, associated with a “token bucket” of size σ tokens. Whenever the bucket is not full, tokens are generated at a rate ρ tokens per second, with $\rho < C$. The link can transmit an arriving packet of size L bits only if the token bucket has at least L tokens - upon the transmission of the packet, the shaper consumes L tokens from the bucket. So, if we start with a full token bucket of size σ tokens, and with a large burst of packets of size L bits each (suppose that σ is an integer multiple of L for simplicity), the link will be able to transmit k of those packets at the rate of the capacity C , with $k = \frac{\sigma/L}{1-\rho/C}$. After those k packets, the link will start transmitting packets at the token generation rate ρ . Usually ρ is referred to as the “shaping rate”, the capacity C is also referred to as the “peak rate”, while σ is referred to as the “maximum burst size”. Another way to describe a traffic shaper is by specifying that the *maximum number of bits* that can be transmitted in any interval of duration τ , starting with a full token bucket, is:

$$\hat{A}(\tau) = \min\{L + C\tau, \sigma + \rho\tau\}$$

The difference between a traffic shaper and a traffic policer is that the former has a buffer to hold packets that arrive when the token bucket is empty [6, 23]. A policer simply drops such “non-conforming” packets. In other words, a shaper delays packets that exceed the traffic shaping profile (σ, ρ) , while a policer drops them.² Policers can cause excessive packet losses and so shapers are more common in practice - we focus on the latter in the rest of the paper.

Why would a residential ISP deploy traffic shaping? First, to allow a user to exceed the service rate that he/she has paid for, for a limited burst size. In that case the user pays for ρ bps, with the additional service capacity $C - \rho$ marketed as a free service enhancement. This is, for instance, how Comcast advertises their PowerBoost traffic shaping mechanism [5]. Second, an ISP may want to limit the service rate provided to the aggregate traffic produced or consumed by a customer, or to limit the service rate consumed by a certain application (e.g. BitTorrent). This form of shaping is relevant to the *network neutrality* debate. Third, certain ISPs prefer to describe their service rates as upper bounds

¹When it is not important to distinguish between shaping and policing, we will simply refer to such mechanisms as “traffic shapers” or just “shapers”.

²A shaper also drops packets once its droptail buffer is full.

for what the user will actually get, e.g., a downstream rate of *at most* 6Mbps. In that case, a shaper can be used to enforce the upper bound of the service rate.

The contribution of this paper is twofold. First, we develop an *active end-to-end detection mechanism*, referred to as *ShaperProbe*, that can infer whether a particular path is subject to traffic shaping, and in that case, estimate the shaper characteristics C , ρ and σ . Second, we analyze results from a large-scale deployment of ShaperProbe on M-Lab [11] since May 2009, detecting traffic shaping in several major ISPs. Our deployment received about one million runs over the last two years from more than 5,700 ISPs; we currently see 2,000-3,000 runs per day (see Figure 1). All data collected through ShaperProbe runs is publicly available through M-Lab [8].³

Traffic shaping detection and estimation methods can be used in different ways: *as a library (API)*; and as a *service* that enables users/administrators to detect or verify their SLAs/shaping configurations. In this paper, we focus on the latter. The ShaperProbe client is a download-and-click userspace binary (no superuser privileges or installation needed) for 32/64-bit Windows, Linux, and OS X; a plugin is also available for the *Vuze* BitTorrent client. The non-UI logic is about 6000 lines of open source native code.

There are several challenges that one needs to tackle when designing an active measurement service that can scale to thousands of users per day, including *accuracy*, *usability* and *non-intrusiveness*. Even though these challenges are often viewed as not significant, at least from the research perspective, they have greatly influenced several design choices and parameter values in ShaperProbe.

Related work. Lakshminarayanan *et al.* recorded initial observations of traffic shaping in residential ISPs in 2003 [18]. Dischinger *et al.* found downstream traffic shaping in a 2007 study [13] using a 10s flow at 10Mbps; they did not find evidence of upstream traffic shaping. More recently, two studies [12, 21] focused on the effects of PowerBoost in cable ISPs. Recent research efforts on detecting traffic discrimination in ISPs [14, 16, 19, 20, 22, 24, 25] complement our work, since we consider the problem of detecting traffic shaping independent of differentiation. In particular, Glasnost [14] uses aggregate throughput of an emulated application and compares it with a baseline to *detect* throughput differences. Weinsberg *et al.* infer queue weights of a discriminatory scheduler [24].

In the rest of the paper, we describe the active detection method (§2), implementation and deployment of ShaperProbe (§3), and look at ShaperProbe data using case studies of four ISPs (§4).

2. ACTIVE PROBING METHOD

The active probing method is an end-to-end process in which one end-host $SN\mathcal{D}$ sends packets on the network path to the receiver \mathcal{RCV} . We detect the presence of traffic shaping in the path $SN\mathcal{D} \rightarrow \mathcal{RCV}$ at \mathcal{RCV} .

Suppose that the narrow link’s capacity on the path is C , and that the sender probes at a *constant bit rate* $R_s = C$. The ShaperProbe capacity estimation process is described in

³We log per-packet send and receive timestamps and sequence numbers for all probing phases, and client IP address and server timestamp (UTC) for each run.

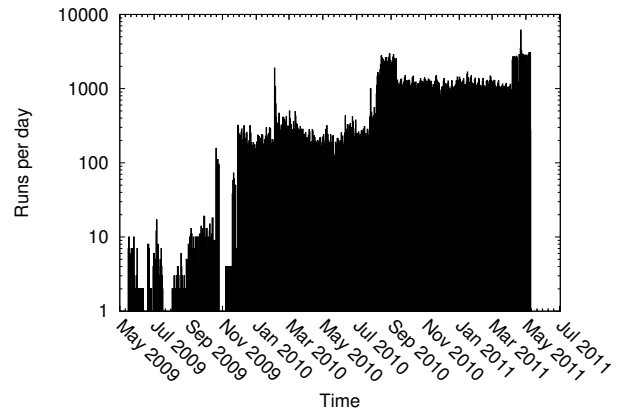


Figure 1: ShaperProbe: volume of runs. The gaps in time show downtime due to tool enhancements.

the Tech Report [17]. The receiver \mathcal{RCV} records the received rate timeseries $R_r(t)$. We compute $R_r(t)$ by discretizing time into fixed size non-overlapping intervals of size Δ . For simplicity, assume that the probing starts at $t = 0$, and that intervals are numbered as integers $i \geq 1$. The i ’th interval includes all packets received in the interval $[(i - 1)\Delta, i\Delta)$, where packet timestamps are taken at \mathcal{RCV} upon receipt of each packet. The discretized received rate timeseries $R_r(i)$ is estimated as the total bytes received in interval i divided by Δ . Note that this estimator of $R_r(t)$ can result in an error of up to $\epsilon = \pm S/\Delta$ where S is the MTU packet size. By choosing a reasonably large Δ , we can reduce the magnitude of ϵ relative to the true received rate.

In the presence of a token bucket traffic shaper (or policer) on $SN\mathcal{D} \rightarrow \mathcal{RCV}$, there exists a value of $i > 1$ at which the received rate timeseries $R_r(i)$ undergoes a *level shift to a lower value*. Our goal is to detect the presence of a level shift, and estimate the token bucket parameters using $R_r(i)$.

2.1 Detection

We want to detect a level shift in R_r in *real-time*, i.e., as we compute the received rate for each new interval. Note that the receiver \mathcal{RCV} is also receiving new packets during the level-shift detection process, and so our method should be fast and computationally light-weight to avoid the introduction of timestamping jitter. The detection method is rather simple and relies on nonparametric rank statistics of R_r so that it is robust to outliers [15].

We compute ranks online. Suppose that we have estimated n values of R_r so far. At the start of the new interval $n + 1$ (i.e., after the receipt of the first packet in that interval), we compute $R_r(n)$ and update the ranks $r(i)$ of $R_r(i)$ for $i = 1 \dots n$. We identify τ as the *start of level shift* if it is the first interval index that satisfies the following three conditions.

First, all ranks at the left of τ are *equal to or higher* than all ranks at the right of τ :

$$\min_{i=1 \dots \tau-1} r(i) \geq \max_{j=\tau+1 \dots n} r(j) \quad (1)$$

Second, we have observed a minimum time duration before and after the current rate measurement:

$$n_L < \tau < n - n_R \quad (2)$$

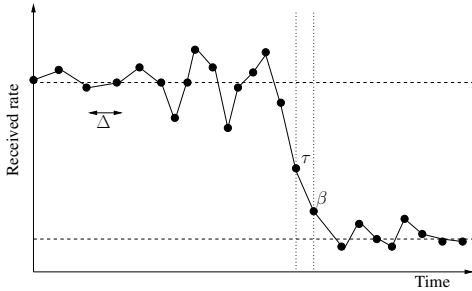


Figure 2: Active probing: Level shift detection.

The value of n_L is chosen based on empirical observations of typical burst durations in ISP deployments, and n_R is a sanity check to ensure that the drop in rate is not just a temporary variation (e.g., due to cross traffic).

Third, we require that there is a *significant drop* in the median rate at point τ :

$$\tilde{R}_r(i) > \gamma \tilde{R}_r(j) \quad (3)$$

$i=1\dots\tau$ $j=\tau\dots n$

where \tilde{R}_r denotes the median, and γ is a suitable threshold. We select γ based on empirical observations of ISP capacities and shaping rates in practice (see Section 2.3).

Similarly, we detect the *end of a level shift* index β such that $\beta \geq \tau$ and β is the *last* point which satisfies the rate condition in Equation 1. Figure 2 illustrates the two level shift indices.

2.2 Estimation

After the detection of a level shift, we estimate the token bucket parameters from the rate timeseries R_r as follows. The token generation rate (shaping rate) ρ is estimated as the *median* (to be robust to outliers) of the received rate measurements *after* β :

$$\hat{\rho} = \tilde{R}_r(i) \quad (4)$$

$i=\beta+1\dots n$

We estimate the token bucket depth (burst size) σ based on the number of bytes sent till the τ 'th time interval. We estimate a range for σ , since we discretize time in to intervals of size Δ , based on the estimate $\hat{\rho}$ of ρ and the received rates:

$$\hat{\sigma} = \sum_{i=1}^{\tau} [R(i) - \hat{\rho}] \Delta \pm \frac{[R(i) - \hat{\rho}] \Delta}{2} \quad (5)$$

2.3 Parameter Selection

As in any other measurement tool that is used in practice, there are some parameters that need to be tuned empirically. In ShaperProbe, the key parameters are the factor γ , the probing duration Λ , and the interval duration Δ . We have selected the values of these parameters based on the detection of actual shaper deployments in broadband ISPs for which we knew the “ground truth”.

Figure 3 shows the ratio of the capacity over the shaping rate C/ρ and the maximum burst duration (in seconds) for 36 advertised traffic shaper deployments at Comcast and Cox in metro Atlanta in October 2010. Note that all tiers have a capacity-to-shaping rate ratio of 1.1 or higher; in the current implementation of ShaperProbe we use $\gamma = 1.1$.

The probing duration Λ should be sufficiently long so that it can detect as many ISP shaping configurations as possible,

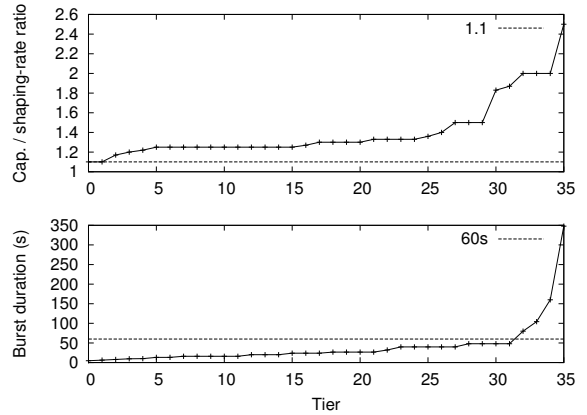


Figure 3: Advertised Comcast, Cox tiers: required γ and Λ .

while at the same time keeping the total probing duration reasonably short when there is no shaping. Figure 3 shows that the burst duration is at most 48s, except for 4 out of 36 configurations. Λ is set to 60s in the current implementation. A typical ShaperProbe residential run lasts for 2-3 minutes.

The averaging window size Δ should be sufficiently large to keep the estimation noise in R_r low, and sufficiently short so that Λ includes several rate samples. We have performed 100 trials in the upstream direction of a Comcast residential connection, whose SLA we know (4.5Mbps shaped to 2Mbps). We found that for $\Delta \geq 50$ ms, the shaping detection rate is 100%; as Δ approaches the inter-packet gap, the detection rate drops significantly. We set Δ to 300ms so that we can detect shaping even in low capacity links.

3. SHAPERPROBE IMPLEMENTATION

The design of a tool that works well on a wide variety of network conditions, OS platforms and broadband link technologies is challenging. A first challenge is that ShaperProbe requires a fast and accurate estimate of the narrow-link capacity between the sender and receiver; this estimate is the ShaperProbe probing rate. ShaperProbe uses packet train dispersion for estimating capacity; it additionally probes using a longer train to be robust to wireless link effects. Second, the probing method should be able to generate traffic at a *constant rate*, even with a coarse-grained userspace OS timer granularity. At the same time, the transmission of packets should not impose heavy load on the CPU resources at the sender. ShaperProbe sends small periodic packet trains, and times the inter-train gaps such that busy-wait loops are minimized. Third, the ShaperProbe client should be non-intrusive. The client and server abort the probing process if they observe losses on the path. Finally, cross traffic on the path may lead to temporary drops in the received rate R_r ; we need to incorporate a filtering mechanism that can remove *outliers* from R_r . ShaperProbe filters outliers using recorded observations from the local neighborhood of the R_r timeseries. The Tech Report [17] describes how ShaperProbe addresses the previous challenges and implementation details.

We currently run load-balanced ShaperProbe server replicas on 48 M-Lab hosts connected directly to tier-1 ASes. For measurement accuracy, we allow only one client at each server replica at any time.

ISP	Upstream (%)	Dwnstrm. (%)
Comcast	71.5 (34874)	73.5 (28272)
Road Runner	6.5 (7923)	63.9 (5870)
AT&T	10.1 (8808)	10.9 (7748)
Cox	63 (5797)	47.4 (4357)
MCI-Verizon	5.6 (8753)	8.4 (7733)

Table 1: Shaping detections: top-5 ISPs in terms of ShaperProbe runs. For each ISP we show percentage of runs with detected shaping and number of total runs.

4. RESULTS

In this section, we take a first look at results from the ShaperProbe deployment at M-Lab. We first examine accuracy using two ISPs for which we know the shaping ground truth and from emulation experiments.

Accuracy. We test the latest version of ShaperProbe on two residential ISPs, AT&T and Comcast, at two homes in metro Atlanta. We use the High-Speed Internet service of Comcast, and the DSL service of AT&T. At the time of these experiments, the Comcast configuration was: {10Mbps up, 22Mbps down} shaped to {2Mbps up, 12Mbps down} [5], while the AT&T configuration did not use shaping ({512Kbps up, 6Mbps down}) [2]. Out of 60 runs, we did not observe any shaping detection errors in either direction at the AT&T connection, while we observed two upstream false negatives at the Comcast connection due to capacity underestimation.

We also emulated token bucket shaping on a wide-area path between a residential Comcast connection and a server deployed at the Georgia Tech campus. We use the LARTC tc tool on Linux with a 2.6.22 kernel on a dual-NIC 1GHz Celeron router with 256MB RAM. Over 20 experiments for each token bucket configuration and 10 configurations, we found that ShaperProbe detects the traffic shaper in all (200) experiments; it also accurately estimates the shaping rate and bucket depth for all configurations.

Data preprocessing. In the following, we analyze data collected from the ShaperProbe M-Lab service. First, we consider runs from the latest ShaperProbe release, collected between 20th October 2009 and 9th May 2011 (total of 845,223 runs). Each run’s trace contains per-packet timestamps and sequence numbers for the upstream and downstream probing “half runs”. Second, we say that a half run is “unfinished” if no shaping was detected and the run lasted for less than 50s - we discard such runs. All completed half runs which are not diagnosed as shaping are considered *no-shaping* cases. Recall that ShaperProbe probes each direction for 60s, and terminates a half run if it either detected shaping or if it observed packet losses during probing. A half run can also be unfinished if the user aborted the client before it could run to completion. After preprocessing, we have a total of 281,394 upstream and 236,423 downstream finished half runs.

Next, we cluster AS numbers into ISPs using their **whois** AS names. The AS information was obtained from Cymru’s whois database in May 2011. Runs which passed the preprocessing checks come from 5,167 distinct ISPs. The top

C (Mbps)	ρ (Mbps)	σ (MB)	Burst duration (s)
3.5	1	5	16.7
4.8	2	5, 10	15.2, 30.5
8.8	5.5	10	25.8
14.5	10	10	18.8

(a) Upstream.

C (Mbps)	ρ (Mbps)	σ (MB)	Burst duration (s)
19.4	6.4	10	6.4
21.1	12.8	10	10.1
28.2	17	20	14.9
34.4	23.4	20	15.3

(b) Downstream.

Table 2: Comcast: detected shaping properties.

five ISPs in terms of the number of runs as well as the fraction of shaping detections are shown in Table 1.

It should be noted that there are several factors that influence the fraction of shaping detections in an ISP. First, ISPs provide multiple tiers of service; some tiers may not use shaping, while service tiers change frequently. Second, an ISP may not deploy shaping in all geographic markets. Third, the access link type can be a factor: a DSL provider can dynamically change the link capacity instead of doing shaping, while a cable provider is more likely to use shaping since DOCSIS provides fixed access capacities. Fourth, for a given connection, the shaping parameters can be dynamically adjusted based on time or load conditions in the ISP. Fifth, an ISP A can originate the BGP prefixes of a smaller ISP B that deploys shaping (while A does not) - we cannot distinguish A from B based on BGP prefix-to-ASN mapping. We study some of these factors in ISP case studies next. Some ISPs disclose their traffic shaping configurations; in such cases, we can validate our observations.

4.1 Case Study: Comcast

Comcast offers Internet connectivity to homes [5] and enterprises [3], and uses two types of access technologies: cable (DOCSIS 3.0) and Ethernet. In each access category, it offers multiple tiers of service. Comcast shapes traffic using the PowerBoost technology [4].

Shaping profiles. We observed many shaping configurations at Comcast between October 2009 and May 2011. Figure 4 shows the shaping configuration of each run (ordered by capacity). For each run, designated by an “ID”, we plot two points in the top panel for the capacity and the shaping rate; and a point in the bottom panel for the burst size. The capacities form an *envelope* of the shaping rate points. We see that there are strong *modes* in the data; Table 2 is a summary of these modes. For higher capacities, we see a larger number of modes in the shaping rate. However, at the tail of the capacity distribution there is only one shaping rate that corresponds to the highest service tier provided by Comcast. We verified our observations with the Comcast website listings [3, 5]. Note that we may not observe all service tiers in that web page, depending on the number of ShaperProbe users at each service tier. We also observe two or three burst sizes that are used across all tiers; the PowerBoost FAQ mentions 10MB and 5MB burst sizes [4].

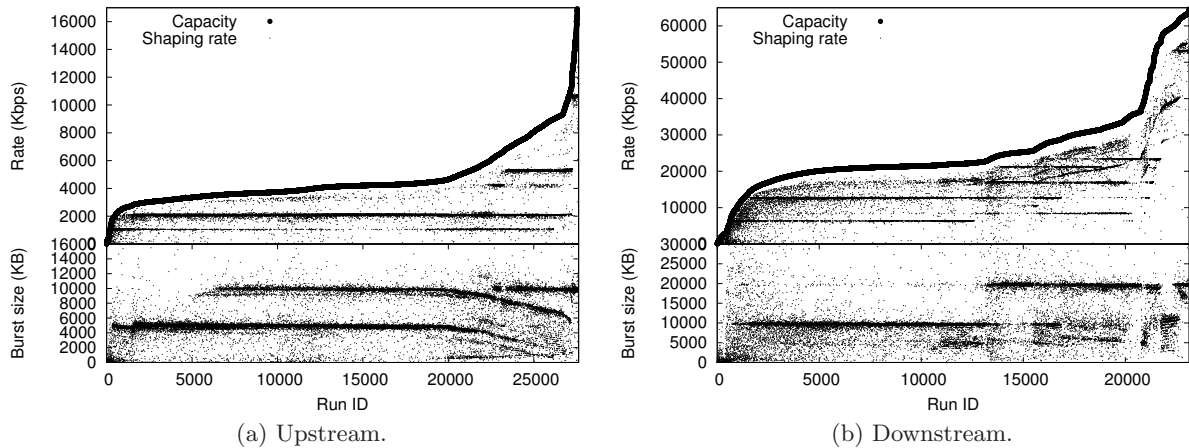


Figure 4: Comcast: Shaping characteristics.

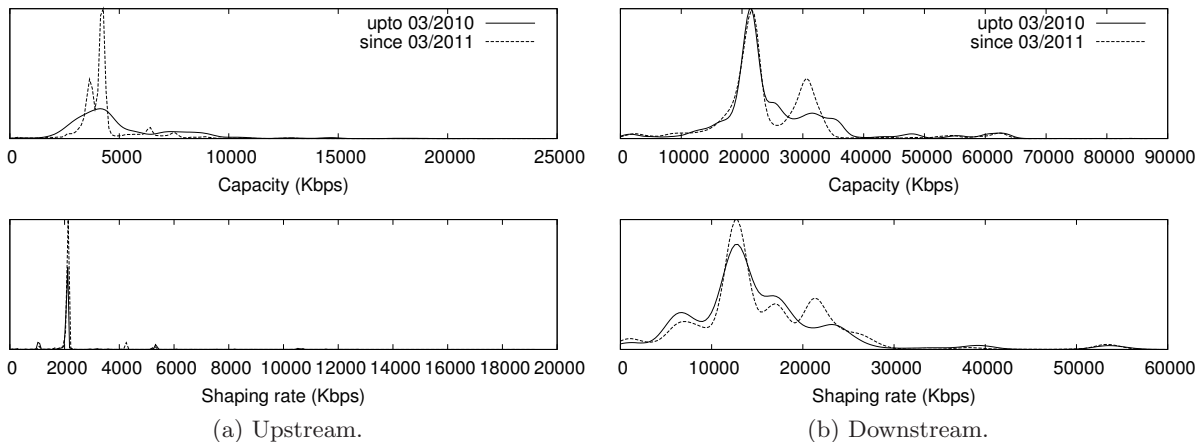


Figure 5: Comcast: histogram of bandwidth with time.

Note that the capacity curves do not show strong modes, unlike the shaping rates. This is due to the underlying DOCSIS access technology. The cable modem uplink is a non-FIFO scheduler; depending on the activity of other nodes at the CMTS, the capacity can vary due to customer scheduling and DOCSIS concatenation. A DOCSIS downlink can also influence the dispersion-based capacity estimates under heavy traffic load conditions because it is a broadcast link.

Did shaping configurations change during the last two years?. We compare data from Comcast collected in October 2009-March 2010 and in March-May 2011. Figure 5 shows estimates of the capacity and shaping rate distributions using a Gaussian kernel density estimator. In the upstream direction, the capacity and shaping rates (the modes of the corresponding distributions) have not changed significantly. The downstream links show a new capacity mode of 30Mbps and a shaping rate mode of 22Mbps in 2011. We did not find significant changes in the burst size during the last two years.

Non-shaped runs. We examine runs in which ShaperProbe did not detect shaping. Figure 6 compares the capacity distribution in such runs with the shaping rate distribution in shaping runs. The non-shaped capacity distributions are *similar* to the shaping rate distributions. Non-shaping runs occur due to the following two reasons. First, Comcast provides service tiers that do not include PowerBoost, but have capacities similar to PowerBoost service tiers (e.g., the Ethernet 1Mbps and 10Mbps business service). Second, it is possible that cross traffic resulted in an empty token bucket at the start of the measurement, and so the capacity that ShaperProbe estimated was equal to the shaping rate; we would not detect shaping in that case.

4.2 Case Studies: Road Runner and Cox

Road Runner (RR) is a cable ISP. A unique aspect of RR is that we have found evidence of downstream shaping, but *no evidence* of upstream shaping in any service tier on their web pages. The ShaperProbe measurements for RR support this observation - 94% of the upstream runs *did not* detect shaping, while 64% of the downstream runs did. Another interesting aspect of RR is that shaping depends on

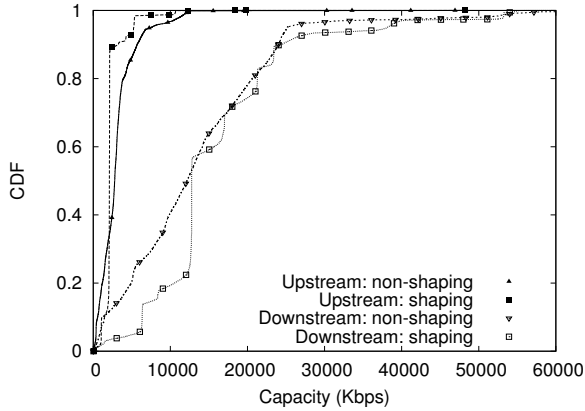


Figure 6: Comcast: distribution of capacities in non-shaping and shaping rate in shaping runs.

the geographic region of the customer; for example, in Texas, RR provides four service tiers: the lower two are not shaped while the upper two are shaped [10]. Under the hypothesis that RR does not shape upstream traffic, we can say that our false positive detection rate for their upstream links is about 6.4%. The capacity distribution of non-shaped RR runs shows that, unlike Comcast, the downstream capacity mode of 750Kbps is not equal to any of their shaping modes (figure omitted due to space constraints; see [17]).

Cox provides residential and business Internet access using cable and Ethernet access. The website [7] mentions that the residential shaping rates and capacities depend on the location of the customer. We gathered residential shaping configurations from the residential pages [7]. The upstream shaping properties of Cox runs in Figure 7 agree with some of the ground truth information we found: (C, ρ) Mbps: (1, 0.77), (1.3, 1), (2, 1.5), (2.5, 1), (2.5, 2), (3, 2), (3.5, 3), (5, 4) and (5.5, 5). Note that the previous ground truth was collected in October 2010, while the ShaperProbe data covers two years. We also found a single burst size mode.

4.3 Case Study: AT&T

Our final case study is that of an ISP for which we do not see frequent shaping detections (10% or less). AT&T provides Internet access to a wide range of customers, from homes and small businesses to enterprises (including other ISPs). Their residential service includes four DSL service tiers [1, 2]. We did not find any mention of traffic shaping in the AT&T service descriptions [1, 2].

Capacity. We first look at the 90% of the runs that did not see shaping. The capacity distribution of non-shaped runs is shown in Figure 8. Given the point-to-point nature of DSL links, ShaperProbe estimates the narrow link capacity more accurately than in cable links. The capacity distributions show several modes: {330Kbps, 650Kbps, 1Mbps, 1.5Mbps} upstream, and {1Mbps, 2.5Mbps, 5Mbps, 6Mbps, 11Mbps, 18Mbps} downstream. We did not observe significant changes in the capacity modes between 2009-2011.

Shaping runs. We look at the 10% of AT&T runs that were *probably mis-diagnosed* as shaping. We found that about a third of these runs exhibit strong shaping rate modes

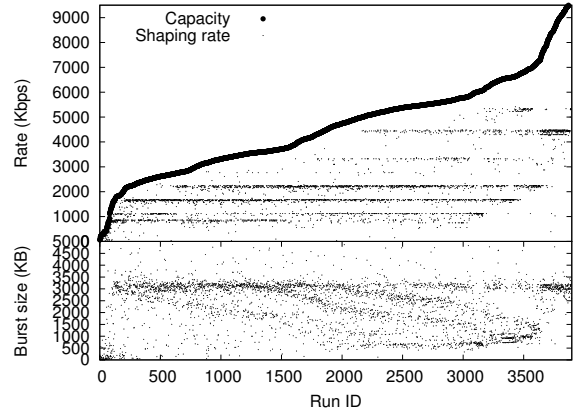


Figure 7: Cox: Upstream shaping.

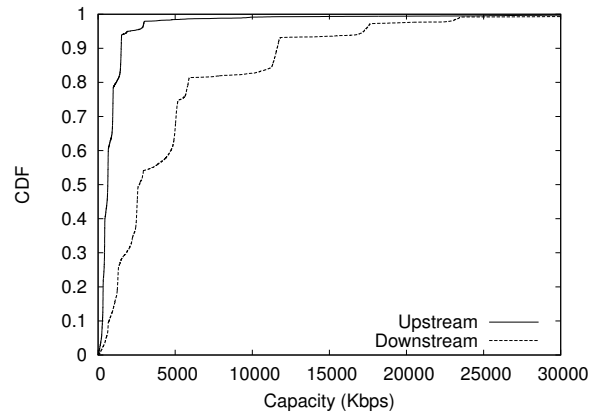


Figure 8: AT&T: Capacity of non-shaping runs.

and an associated burst size mode (for figure, see [17]). About 80% of 333 runs that have the shaping rate modes come from hostnames that resolve to the domain `mchsi.com`, owned by the cable ISP Mediacom [9]. So, it is possible (though we can not be certain) that these shaping detections were *not* errors afterall.

5. CONCLUSION

In this work, we presented an end-to-end active probing, detection, and estimation method of traffic shaping in ISPs. Our evaluation using controlled experiments and in two known ISP deployments shows that ShaperProbe has false positive and false negative detection rates of less than 5%. We presented a first large-scale study of shaping at four large ISPs, and validated some of our observations using ISP advertised tier data. A strong modality of shaping rates and burst sizes suggests that ISPs typically deploy a small set of shaping configurations. We found some shaping detections for which the ISPs do not mention shaping in their service descriptions⁴. Lack of publicly-available information, however, does not necessary imply that these are false detections. We are currently working on passive shaping detection methods (for preliminary results, see [17]).

⁴ISPs, however, typically mention in their SLAs that “listed capacities may vary”.

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